# Construction the Linear Codes in Projective Plane of Order Sixteen

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Article Info	Abstract
Page Number: 2070 - 2087 Publication Issue:	The major aim of this research is to introduce the maximum value of size of complete (n;r)-arc to be existence in the projective plane of order
Vol 71 No. 4 (2022)	sixteen $PG(2,16)$ and then study the relationship between coding theory
	and a finite projective plane, so apply the results of complete $(n; r)$ -arc;
Article History	$r = 4, \dots, 17$ to coding theory.
Article Received: 25 March 2022	Keywords: Projective Plane and coding theory .
Revised: 30 April 2022	
Accepted: 15 June 2022	
Publication: 19 August 2022	

**1.Introduction.** The aim of coding theory is to develop methods that enable the recipient of message to detect or even correct that occur while transmitting data. Many aspects of coding theory can be directly translated into geometry problems. A linear  $[n; k, n - r]_q$ -code is an n-dimensional subspace of the k-dimensional vector space V(k,q) with non zero vectors weight at least d. An important problem in coding theory is that to optimise one of the parameters k, n, d for given value of the other two and fixed q. The subject of arcs is not only interesting in its purely geometrical setting. An (n;r)-arcs have applications in coding theory, where they can be interpreted as a linear  $[n; n - r]_q$ -code. So every  $[n; 3, n - r]_q$ -code is equivalent to (n; r) - arc in PG(2, q) containing at least r collinear points. In [5] R. Hill studied a the fundamentals of coding theory. In [7] Al-Zangana has been studied the geometry of the plane of order nineteen and its application to error -correcting codes. For more details see [11],[3]. The aim of this research is to construct the complete (n; r) – arcs

where r = 4, ..., 17, and construct the all codes corresponding to these arcs. All calculations are by Gap program[4].

#### 2. The projective plane PG(2,16).

In PG(2,16) there are 273 points and lines, 17 points on each line and 17 lines passage through each point. Take  $H(X) = X^2 + X + \omega^2$  in  $F_{16}[X]$ , where

$$F_{16} = \{0, 1, \omega, \omega^2, \omega^3, \omega^4, \omega^5, \omega^6, \omega^7, \omega^8, \omega^9, \omega^{10}, \omega^{11}, \omega^{12}, \omega^{13}, \omega^{14}\}$$

a polynomial is primitive in  $F_{16}$ . The companion matrix of H is  $T = \begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ \omega^7 & 1 & 0 \end{bmatrix}$ . The points of PG(2,16) are generated by T as follows:  $P_i = P[1,0,0]T^i$ ;  $i = 0, \dots 272$ .

To find the lines in PG(2,16) :Let  $l_0$  contains of 17 points such that the third coordinate of it

is equal to zero. Then the points  $P_i = i$  and the lines  $l_i$  in PG(2,16) can be represented by:

 $l_{0}=\{0,1,3,7,15,31,63,90,116,127,136,181,194,204,233,138,255\}.$  Moreover,  $l_{i}=l_{0}T^{i}$ ;  $i=0,\ldots 272$ .

#### 3.Some definitions and basic properties

**Definition 3.1.**[6] An (n; r)-arc  $\mathcal{K}$ in PG(2,q) is a set of n points, satisfies that every line meet it in less than or equal r points, that is  $|K \cap l| \le r$  for all  $l \in PG(2,16)$ .

**Definition 3.2**[9]:The points out of arc K which passes through it *i* bisecant of K is called a point of index *i*.The number of these points is denoted by  $c_i$ .So which represents the number of the points not on bisecant of K.

**Lemma 3.1:** For a (n, r) - arc K, the following equations hold:

$$\sum_{i=0}^{n'} c_i = q^2 + q + 1 - n \tag{1.1}$$

$$\sum_{i=0}^{n} ic_i = K(k-1)(q-1)/2 \tag{1.2}$$

$$\sum_{i=0}^{n} i(i-1)c_i = K(K-1)(K-2)(K-3)/8$$
(13)

Vol. 71 No. 4 (2022) http://philstat.org.ph 2071

such that 
$$n = \left\lfloor \frac{1}{2}n \right\rfloor$$

Proof. See[1].

**Definition 3.3.[1]** An linear  $[n, k, d]_q$  -code *C* over a finite field is subspace of dimension *k* of the n – dimensional vector space  $V(n, q) = F_q^n$  such that any two distinct vectors in *C* differ in at least of *d* places. The elements of the code are called codewords. Also the parameters *n*, *k* and *d* are called length, dimension, and minimum distance of *C*.

**Definition 3.4.[5]** For any two code words the minimum distance (Hamming distance) between  $c_1$  and,  $c_2$  is denoted by  $d(c_1, c_2)$  and it is defined to be the number of positions in which the corresponding coordinates differ. The minimum distance of C is  $d(C) = \min\{d(c_1, c_2); c_1, c_2 \in C, c_1 \neq c_2\}$ .

**Definition 3.5.[5]** The weight w(x) of  $x \in V(n,q)$  is w(x) = d(x,0); that is, w(x) is the number of non-zero elements in x.

Lemma 3.1.  $d(c_1, c_2) = w(c_1 - c_2)$  for  $c_1, c_2 \in C$ .

Proof. See [1].

**Definition 3.6.[5]** Let *C* be a linear  $[n, k, d]_q$  -code and  $A_i$  be the number of codewords of weight *i* in a code *C*, the list  $A_i$  for  $0 \le i \le n$  is called the weight distribution of *C*.

**Definition 3.7.[1]** Two linear codes  $C_1$  and  $C_2$  in V(n,q) are equivalent if  $C_1$  can be obtained from  $C_2$  by permuting coordinates and by multiplying coordinates by non-zero elements of  $F_q$ .

**Definition 3.8.[1]** A generator matrix of a linear  $[n, k, d]_q$ -code *C* is  $k \times n$  matrix over the finite field  $F_q$  whose rows from a basis of *C*; it is denoted by *G*.

**Definition 3.9.[1]** For a linear  $[n, k, d]_q$  -code *C* over the finite field  $F_q$ , the singleton bound that  $d(C) \le n - k + 1$ .

**Definition 3.10.[1]** A linear[n, k, d]<sub>q</sub> -code C over a finite field is said to be maximum distance separable (MDS) code if d satisfies the following bound :

$$d(C) = n - k + 1.$$

And if d = n - k, then the code is called almost maximum distance separable (AMDS).

**Theorem 3.1.[5]** There exists a projective  $[n, K, d]_q$ -code if and only if there exists a (n, n - d)-arc.

**Definition 3.11**.[1] For  $x_0 \in F_q^n$  and  $r \in Z, r \ge 0$ , the ball of centre  $x_0$  and radius r is  $S(x_0, r) = \{x \in F_q^n; d(x_0, x) \le r\}$ 

**Definition 3.12.[1]** The covering radius of linear  $[n, k, d]_q$  code *C* is the smallest  $\mu = \mu(C)$  such that  $\bigcup_{x \in C} S(x, \mu) = F_q^n$ .

#### 4. The construction of complete arc of higher degree in projective plane of order sixteen

**Theorem 4.1.** In projective plane PG (2,16) there exists :

I. A complete (36;4)-arc.

II. A complete (46;5)-arc

III. A complete (57;6)-arc

IV. A complete (72;7)-arc

V. A complete (83;8)-arc

VI. A complete (100;9)-arc

VII. A complete (118;10)-arc

VIII. A complete (132;11)-arc

IX. A complete (154;12)-arc

**X.** A complete (175;13)-arc

**XI.** A complete (192,14)-arc

**XII.** A complete (213;15)-arc

XIII. A complete (234;16)-arc

XIV. A complete (273;17)-arc

## **Proof:**

I. we choose the (24;3) -arc and then intersect it with the lines in PG(2,16) when r = 4.

The number of the points of  $c_i$  which is represents the number of points out of arc K which are passes through them *i* 2 -secants of (24,4)-arc as follows:

 $[c_0, \dots, c_{12}] = [249, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0]$ 

Then by adding the first point of the  $c_0$  of (24;4)-arc to (24;4)-arc to construct (25;4)-arc, the values of parameters  $c_i$  of (25;4)-arc ; i = 0, ..., 12 are respectively,  $[c_0, ..., c_{12}] = [183,65,0,0,0,0,0,0,0,0,0,0,0]$ ;

Then by adding the first point of the  $c_0$  of (25;4)-arc to (25;4)-arc to construct (26;4)-arc, the values of parameters  $c_i$  of (26;4)-arc ; i = 0, ... 13 are respectively,  $[c_0, ..., c_{13}] = [180,65,0,0,0,0,0,0,0,0,0,0,0]$ ;

Then continue in the same way until we get (36;4)-arc which is complete since  $c_0 = 0$ .

**II.** we choose the (24;3) -arc and then intersect it with the lines in PG(2,16) when r = 5.

The number of the points of  $c_i$  which is represents the number of points out of arc K which are passes through them *i* 2 -secants of (24,5)-arc as follows:

 $[c_0, \dots, c_{12}] = [249, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0]$ 

Then by adding the first point of the  $c_0$  of (24;5)-arc to (24;5)-arc to construct (25;5)-arc, the values of parameters  $c_i$  of (25;5)-arc; i = 0, ..., 12 are respectively,  $[c_0, ..., c_{12}] = [148,0,0,0,0,0,0,0,0,0,0,0,0,0];$ 

Then by adding the first point of the  $c_0$  of (25;5)-arc to (25;5)-arc to construct (26;5)-arc, the values of parameters  $c_i$  of (26;5)-arc; i = 0, ..., 13 are respectively,  $[c_0, ..., c_{13}] = [146,0,0,0,0,0,0,0,0,0,0,0,0,0];$ 

Then continue in the same way until we get (46;5)-arc which is complete since  $c_0 = 0$ .

**III.** we choose the (24;3) -arc and then intersect it with the lines in PG(2,16) when r = 6.

The number of the points of  $c_i$  which is represents the number of points out of arc K which are passes through them *i* 2 -secants of (24,6)-arc as follows:

 $[c_0, \dots, c_{12}] = [249, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0]$ 

Then by adding the first point of the  $c_0$  of (24;6)-arc to (24;6)-arc to construct (25;6)-arc, the values of parameters  $c_i$  of (25;6)-arc ; i = 0, ... 12 are respectively,  $[c_0, ..., c_{12}] = [148,0,0,0,0,0,0,0,0,0,0,0,0]$ ;

Then by adding the first point of the  $c_0$  of (25;6)-arc to (25;6)-arc to construct (26;6)-arc, the values of parameters  $c_i$  of (26;6)-arc; i = 0, ... 13 are respectively,  $[c_0, ..., c_{13}] = [146,0,0,0,0,0,0,0,0,0,0,0,0,0];$ 

Then continue in the same way until we get (57;6)-arc which is complete since  $c_0 = 0$ .

V. we choose the (24;3) -arc and then intersect it with the lines in PG(2,16) when r = 7.

The number of the points of  $c_i$  which is represents the number of points out of arc K which are passes through them *i* 2 -secants of (24,7)-arc as follows:

 $[c_0, \dots, c_{12}] = [249, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0]$ 

Then by adding the first point of the  $c_0$  of (24;7)-arc to (24;7)-arc to construct (25;7)-arc, the values of parameters  $c_i$  of (25;7)-arc ; i = 0, ... 12 are respectively,  $[c_0, ..., c_{12}] = [148,0,0,0,0,0,0,0,0,0,0,0,0]$ ;

Then by adding the first point of the  $c_0$  of (25;7)-arc to (25;7)-arc to construct (26;7)-arc, the values of parameters  $c_i$  of (26;7)-arc; i = 0, ..., 13 are respectively,  $[c_0, ..., c_{13}] = [146,0,0,0,0,0,0,0,0,0,0,0,0];$ 

Then continue in the same way until we get (72;7)-arc which is complete since  $c_0 = 0$ .

**IV.** we choose the (24;3) -arc and then intersect it with the lines in PG(2,16) when r = 8.

The number of the points of  $c_i$  which is represents the number of points out of arc K which are passes through them *i* 2 -secants of (24,8)-arc as follows:

 $[c_0, \dots, c_{12}] = [249, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0]$ 

Then by adding the first point of the  $c_0$  of (24;8)-arc to (24;8)-arc to construct (25;8)-arc, the values of parameters  $c_i$  of (25;8)-arc ; i = 0, ..., 12 are respectively,  $[c_0, ..., c_{12}] = [148,0,0,0,0,0,0,0,0,0,0,0,0]$ ;

Then by adding the first point of the  $c_0$  of (25;8)-arc to (25;8)-arc to construct (26;8)-arc, the values of parameters  $c_i$  of (26;8)-arc ; i = 0, ... 13 are respectively,  $[c_0, ..., c_{13}] = [146,0,0,0,0,0,0,0,0,0,0,0,0]$ ;

Then continue in the same way until we get (83;8)-arc which is complete since  $c_0 = 0$ .

V. we choose the (24;3) -arc and then intersect it with the lines in PG(2,16) when r = 9.

The number of the points of  $c_i$  which is represents the number of points out of arc K which are passes through them *i* 2 -secants of (24,9)-arc as follows:

 $[c_0, \dots, c_{12}] = [249, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0]$ 

Then by adding the first point of the  $c_0$  of (24;9)-arc to (24;9)-arc to construct (25;9)-arc, the values of parameters  $c_i$  of (25;9)-arc ; i = 0, ..., 12 are respectively,  $[c_0, ..., c_{12}] = [148,0,0,0,0,0,0,0,0,0,0,0,0]$ ;

Then by adding the first point of the  $c_0$  of (25;9)-arc to (25;9)-arc to construct (26;9)-arc, the values of parameters  $c_i$  of (26;9)-arc ; i = 0, ..., 13 are respectively,  $[c_0, ..., c_{13}] = [146,0,0,0,0,0,0,0,0,0,0,0,0,0]$ ;

Then continue in the same way until we get (100;9)-arc which is complete since  $c_0 = 0$ .

VI. we choose the (24;3) -arc and then intersect it with the lines in PG(2,16) when r = 10.

The number of the points of  $c_i$  which is represents the number of points out of arc K which are passes through them *i* 2 -secants of (24,10)-arc as follows:

 $[c_0, \dots, c_{12}] = [249, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0]$ 

Then by adding the first point of the  $c_0$  of (24;10)-arc to (24;10)-arc to construct (25;10)-arc, the values of parameters  $c_i$  of (25;10)-arc ; i = 0, ... 12 are respectively,  $[c_0, ..., c_{12}] = [148,0,0,0,0,0,0,0,0,0,0,0,0]$ ;

Then by adding the first point of the  $c_0$  of (25;10)-arc to (25;10)-arc to construct (26;10)-arc, the values of parameters  $c_i$  of (26;10)-arc; i = 0, ... 13 are respectively,  $[c_0, ..., c_{13}] = [146,0,0,0,0,0,0,0,0,0,0,0,0,0]$ ;

Then continue in the same way until we get (118;10)-arc which is complete since  $c_0 = 0$ . VII. we choose the (24;3) -arc and then intersect it with the lines in PG(2,16) when r = 11. The number of the points of  $c_i$  which is represents the number of points out of arc K which are passes through them *i* 2 -secants of (24,11)-arc as follows:

 $[c_0, \dots, c_{12}] = [249, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0]$ 

Then by adding the first point of the  $c_0$  of (24;11)-arc to (24;11)-arc to construct (25;11)-arc, the values of parameters  $c_i$  of (25;11)-arc ; i = 0, ... 12 are respectively,  $[c_0, ..., c_{12}] = [148,0,0,0,0,0,0,0,0,0,0,0,0]$ ;

Then by adding the first point of the  $c_0$  of (25;11)-arc to (25;11)-arc to construct (26;11)-arc, the values of parameters  $c_i$  of (26;11)-arc; i = 0, ... 13 are respectively,  $[c_0, ..., c_{13}] = [146,0,0,0,0,0,0,0,0,0,0,0,0,0]$ ;

Then continue in the same way until we get (132;11)-arc which is complete since  $c_0 = 0$ .

VIII. we choose the (24;3) -arc and then intersect it with the lines in PG(2,16) when r = 12.

The number of the points of  $c_i$  which is represents the number of points out of arc K which are passes through them *i* 2 -secants of (24,12)-arc as follows:

 $[c_0, \dots, c_{12}] = [249, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0]$ 

Then by adding the first point of the  $c_0$  of (24;12)-arc to (24;5)-arc to construct (25;12)-arc, the values of parameters  $c_i$  of (25;12)-arc ; i = 0, ... 12 are respectively,  $[c_0, ..., c_{12}] = [148,0,0,0,0,0,0,0,0,0,0,0,0]$ ;

Then by adding the first point of the  $c_0$  of (25;12)-arc to (25;12)-arc to construct (26;12)-arc, the values of parameters  $c_i$  of (26;12)-arc; i = 0, ... 13 are respectively,  $[c_0, ..., c_{13}] = [146,0,0,0,0,0,0,0,0,0,0,0,0,0];$ 

Then continue in the same way until we get (154;12)-arc which is complete since  $c_0 = 0$ .

**IX.** we choose the (24;3) -arc and then intersect it with the lines in PG(2,16) when r = 13.

The number of the points of  $c_i$  which is represents the number of points out of arc K which are passes through them *i* 2 -secants of (24,13)-arc as follows:

 $[c_0, \dots, c_{12}] = [249, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0]$ 

Then by adding the first point of the  $c_0$  of (24;13)-arc to (24;13)-arc to construct (25;13)-arc, the values of parameters  $c_i$  of (25;13)-arc; i = 0, ... 12 are respectively,  $[c_0, ..., c_{12}] =$ [148,0,0,0,0,0,0,0,0,0,0,0,0];

Then by adding the first point of the  $c_0$  of (25;13)-arc to (25;13)-arc to construct (26;13)-arc, the values of parameters  $c_i$  of (26;13)-arc ; i = 0, ... 13 are respectively,  $[c_0, ..., c_{13}] = [146,0,0,0,0,0,0,0,0,0,0,0,0,0]$ ;

Then continue in the same way until we get (175;13)-arc which is complete since  $c_0 = 0$ .

**X.** we choose the (24;3) -arc and then intersect it with the lines in PG(2,16) when r = 14.

The number of the points of  $c_i$  which is represents the number of points out of arc K which are passes through them *i* 2 -secants of (24,5)-arc as follows:

 $[c_0, \dots, c_{12}] = [249, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0]$ 

Then by adding the first point of the  $c_0$  of (24;14)-arc to (24;14)-arc to construct (25;14)-arc, the values of parameters  $c_i$  of (25;14)-arc ; i = 0, ... 12 are respectively,  $[c_0, ..., c_{12}] = [148,0,0,0,0,0,0,0,0,0,0,0,0]$ ;

Then by adding the first point of the  $c_0$  of (25;14)-arc to (25;14)-arc to construct (26;14)-arc, the values of parameters  $c_i$  of (26;14)-arc; i = 0, ... 13 are respectively,  $[c_0, ..., c_{13}] = [146,0,0,0,0,0,0,0,0,0,0,0,0,0];$ 

Then continue in the same way until we get (192;14)-arc which is complete since  $c_0 = 0$ .

XI. we choose the (24;3) -arc and then intersect it with the lines in PG(2,16) when r = 15.

The number of the points of  $c_i$  which is represents the number of points out of arc K which are passes through them *i* 2 -secants of (24,15)-arc as follows:

 $[c_0, \dots, c_{12}] = [249, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0]$ 

Then by adding the first point of the  $c_0$  of (24;15)-arc to (24;15)-arc to construct (25;15)-arc, the values of parameters  $c_i$  of (25;15)-arc; i = 0, ... 12 are respectively,  $[c_0, ..., c_{12}] = [148,0,0,0,0,0,0,0,0,0,0,0,0]$ ;

Then by adding the first point of the  $c_0$  of (25;15)-arc to (25;15)-arc to construct (26;15)-arc, the values of parameters  $c_i$  of (26;15)-arc; i = 0, ... 13 are respectively,  $[c_0, ..., c_{13}] = [146,0,0,0,0,0,0,0,0,0,0,0,0,0];$ 

Then continue in the same way until we get (214;15)-arc which is complete since  $c_0 = 0$ .

**XII.** we choose the (24;3) -arc and then intersect it with the lines in PG(2,16) when r = 16.

The number of the points of  $c_i$  which is represents the number of points out of arc K which are passes through them *i* 2 -secants of (24,16)-arc as follows:

 $[c_0, \dots, c_{12}] = [249, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0]$ 

Then by adding the first point of the  $c_0$  of (24;16)-arc to (24;16)-arc to construct (25;16)-arc, the values of parameters  $c_i$  of (25;16)-arc ; i = 0, ... 12 are respectively,  $[c_0, ..., c_{12}] = [148,0,0,0,0,0,0,0,0,0,0,0,0]$ ;

Then by adding the first point of the  $c_0$  of (25;16)-arc to (25;5)-arc to construct (26;16)-arc, the values of parameters  $c_i$  of (26;16)-arc ; i = 0, ... 13 are respectively,  $[c_0, ..., c_{13}] = [146,0,0,0,0,0,0,0,0,0,0,0,0,0]$ ;

Then continue in the same way until we get (234;16)-arc which is complete since  $c_0 = 0$ .

**XIII**. we choose the (24;3) -arc and then intersect it with the lines in PG(2,16) when r = 17.

The number of the points of  $c_i$  which is represents the number of points out of arc K which are passes through them *i* 2 -secants of (24,17)-arc as follows:

 $[c_0, \dots, c_{12}] = [249, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0]$ 

Then by adding the first point of the  $c_0$  of (24;17)-arc to (24;17)-arc to construct (25;17)-arc, the values of parameters  $c_i$  of (25;17)-arc; i = 0, ... 12 are respectively,  $[c_0, ..., c_{12}] = [148,0,0,0,0,0,0,0,0,0,0,0,0];$ 

Then by adding the first point of the  $c_0$  of (25;17)-arc to (25;17)-arc to construct (26;17)-arc, the values of parameters  $c_i$  of (26;17)-arc; i = 0, ... 13 are respectively,  $[c_0, ..., c_{13}] = [146,0,0,0,0,0,0,0,0,0,0,0,0,0];$ 

Then continue in the same way until we get (273;17)-arc which is complete since  $c_0 = 0$ .

#### 5. Constructions of the linear codes in projective plane of order sixteen

The problem of determine the largest size of (n; r) –arcs in PG(2,q) demonstrated an interesting connection with coding theory. This connection is between (n; r) –arcs in PG(2,q) and the  $[n, k, d]_q$  codes in coding theory. This link gives the following theorem.

**Theorem 5.1.** In projective plane PG(2,16) there exists:

I. A projective [36,3,32]<sub>16</sub>-code if there exists a (36,4)-arc.

II. A projective [46,3,41]<sub>16</sub>-code if there exists a (41,5)-arc.

III. A projective [57,3,51]<sub>16</sub>-code if there exists a (57,6)-arc.

IV. A projective [72,3,65]<sub>16</sub>-code if there exists a (72,7)-arc.

V. A projective [83,3,75]<sub>16</sub>-code if there exists a (83,8)-arc.

VI. A projective  $[100,3,91]_{16}$ -code if there exists a (100,9)-arc.

VII. A projective [118,3,108]<sub>16</sub>-code if there exists a (118,10)-arc.

VIII. A projective [132,3,121]<sub>16</sub>-code if there exists a (132,11)-arc.

**IX.** A projective  $[154,3,142]_{16}$ -code if there exists a (154,12)-arc.

**X.** A projective [175,3,162]<sub>16</sub>-code if there exists a (175,13)-arc .

XI. A projective [192,3,178]<sub>16</sub>-code if there exists a (192,14)-arc .

XII. A projective [213,3,198]<sub>16</sub>-code if there exists a (213,15)-arc.

XIII A projective [234,3,218]<sub>16</sub>-code if there exists a (234,16)-arc.

XIV. A projective [273,3,256]<sub>16</sub>-code if there exists a (273,17)-arc.

#### Proof.

According to the theorem (3.1) and (4.1) an(n, n - d)-arc in PG(k - 1, q) is equivalent to a projective  $[n, k, d]_q - code$ . Now if q = 16, K = 3 and n - d = r; r is degree of arc, then there is an one to one correspondence between (n; r)- arc in PG(k - 1, q) and a projective  $[n, 3, n - r]_q - code$ .

I. (36;4)-arc give linear  $[36,3,32]_{16}$  – *code* define by the generator matrix  $G_{3\times 36}$ .

$$G_{3\times 36} = \begin{bmatrix} 1 & 0 & 0 & w^7 & 1 & w^{11} & w^{13} & w^4 & w^9 & w^6 \\ 0 & 1 & 0 & 1 & 1 & w^3 & w^7 & w^9 & w^{13} & \dots & w^5 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 1 \end{bmatrix}$$

II. (46;5)-arc give linear  $[46,3,41]_{16}$  – *code* define by the generator matrix  $G_{3\times 46}$ .

$$G_{3\times 46} = \begin{bmatrix} 1 & 0 & 0 & w^7 & 1 & w^{11} & w^{13} & w^4 & w^9 & w^4 & w^{10} & w^2 \\ 0 & 1 & 0 & 1 & 1 & w^3 & w^7 & w^9 & w^{13} & w^9 & w^8 & \dots & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \end{bmatrix}$$

radius  $\mu = 43$ .  $\Box$ 

III. (57;6)-arc give linear  $[57,3,51]_{16}$  – *code* define by the generator matrix  $G_{3\times 57}$ .

$$G_{3\times 57} = \begin{bmatrix} 1 & 0 & 0 & w^7 & 1 & w^{11} & w^{13} & w^4 & w^9 & w^4 & w^{10} & w^6 \\ 0 & 1 & 0 & 1 & 1 & w^3 & w^7 & w^9 & w^{13} & w^9 & w^8 & \dots & w^7 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \end{bmatrix}$$

IV. (72;7)-arc give linear  $[72,3,65]_{16}$  – *code* define by the generator matrix  $G_{3\times72}$ .

 $G_{3\times72} = \begin{bmatrix} 1 & 0 & 0 & w^7 & 1 & w^{11} & w^{13} & w^4 & w^9 & w^4 & w^{10} & w^3 \\ 0 & 1 & 0 & 1 & 1 & w^3 & w^7 & w^9 & w^{13} & w^9 & w^8 & \dots & w \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \end{bmatrix}$ 

Not that  $A_0 = 1$ ,  $A_{65} = 870$ ,  $A_{66} = 465$ ,  $A_{67} = 540$ ,  $A_{68} = 735$ ,  $A_{69} = 930$ ,  $A_{70} = 495$ ,  $A_{71} = 60$ and  $S = A_0 + A_{65} + A_{66} + A_{67} + A_{68} + A_{69} + A_{70} + A_{71} = 4096 = 16^3$  and the covering radius  $\mu = 67$ .  $\Box$ 

V. (83;8)-arc give linear  $[83,3,75]_{16}$  – *code* define by the generator matrix  $G_{3\times83}$ .

 $G_{3\times83} = \begin{bmatrix} 1 & 0 & 0 & w^7 & 1 & w^{11} & w^{13} & w^4 & w^9 & w^4 & w^{10} & w^7 \\ 0 & 1 & 0 & 1 & 1 & w^3 & w^7 & w^9 & w^{13} & w^9 & w^8 & \dots & w^{10} \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \end{bmatrix}$ 

Not that  $A_0 = 1$ ,  $A_{75} = 810$ ,  $A_{76} = 285$ ,  $A_{77} = 510$ ,  $A_{78} = 615$ ,  $A_{79} = 1095$ ,  $A_{80} = 615$ ,  $A_{81} = 165$ and  $S = A_0 + A_{75} + A_{76} + A_{77} + A_{78} + A_{79} + A_{80} + A_{81} = 4096 = 16^3$  and the covering radius  $\mu = 78$ .  $\Box$ 

VI. (100;9)-arc give linear  $[100,3,91]_{16}$  – *code* define by the generator matrix  $G_{3\times 100}$ .

Not that  $A_0=1$ ,  $A_{91}=825$ ,  $A_{92}=480$ ,  $A_{93}=390$ ,  $A_{94}=630$ ,  $A_{95}=900$ ,  $A_{96}=630$ ,  $A_{97}=225$ ,  $A_{98}=15$ , and  $S=A_0+A_{91}+A_{92}+A_{93}+A_{94}+A_{95}+A_{96}+A_{97}+A_{98}=4096=16^3$  and the covering radius  $\mu = 94$ .  $\Box$ 

**VII**. (118;10)-arc give linear  $[118,3,108]_{16}$  – *code* define by the generator matrix  $G_{3\times 118}$ .

$$G_{3\times 118} = \begin{bmatrix} 1 & 0 & 0 \ \text{w}^7 & 1 & \text{w}^{11} \ \text{w}^{13} & \text{w}^4 & \text{w}^9 \ \text{w}^4 \ \text{w}^{10} & \text{w}^5 \\ 0 & 1 & 0 & 1 & 1 & \text{w}^3 \ \text{w}^7 & \text{w}^9 & \text{w}^{13} \ \text{w}^9 \ \text{w}^8 \ \dots \ \text{w}^8 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \end{bmatrix}$$

Not that  $A_0=1$ ,  $A_{108}=915$ ,  $A_{109}=600$ ,  $A_{110}=330$ ,  $A_{111}=540$ ,  $A_{112}=810$ ,  $A_{113}=660$ ,  $A_{114}=240$ , and  $S=A_0+A_{108}+A_{109}+A_{110}+A_{111}+A_{112}+A_{113}+A_{114}=4096=16^3$  and the covering radius  $\mu = 111$ .  $\Box$ 

VIII. (132;11)-arc give linear  $[132,3,121]_{16}$  – *code* define by the generator matrix  $G_{3\times 132}$ .

 $G_{3\times 132} = \begin{bmatrix} 1 & 0 & 0 & w^7 & 1 & w^{11} & w^{13} & w^4 & w^9 & w^4 & w^{10} & 0 \\ 0 & 1 & 0 & 1 & 1 & w^3 & w^7 & w^9 & w^{13} & w^9 & w^8 & \dots & w^{13} \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \end{bmatrix}$ 

Not that  $A_0=1$ ,  $A_{121}=195$ ,  $A_{122}=570$ ,  $A_{123}=660$ ,  $A_{124}=495$ ,  $A_{125}=435$ ,  $A_{126}=690$ ,  $A_{127}=720$ ,  $A_{128}=300$ ,  $A_{129}=30$  and  $S=A_0+A_{121}+A_{122}+A_{123}+A_{124}+A_{125}+A_{126}+A_{127}+A_{128}+A_{128}=4096=16^3$  and the covering radius  $\mu = 124$ .  $\Box$ 

IX. (154;12)-arc give linear  $[154,3,142]_{16}$  – code define by the generator matrix  $G_{3\times 154}$ .

 $G_{3\times 154} = \begin{bmatrix} 1 & 0 & 0 & w^7 & 1 & w^{11} & w^{13} & w^4 & w^9 & w^4 & w^{10} & w^{11} \\ 0 & 1 & 0 & 1 & 1 & w^3 & w^7 & w^9 & w^{13} & w^9 & w^8 & \dots & w^{13} \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \end{bmatrix}$ 

Not that  $A_0=1$ ,  $A_{142}=825$ ,  $A_{143}=1020$ ,  $A_{144}=480$ ,  $A_{145}=345$ ,  $A_{146}=585$ ,  $A_{147}=540$ ,  $A_{148}=285$ ,  $A_{149}=15$ , and  $S=A_0+A_{142}+A_{143}+A_{144}+A_{145}+A_{146}+A_{147}+A_{148}+A_{149}$  $=4096=16^3$  and the covering radius  $\mu = 144$ .  $\Box$ 

**X.** (175;13)-arc give linear  $[175,3,162]_{16}$  – *code* define by the generator matrix  $G_{3\times 175}$ .

$$G_{3\times 175} = \begin{bmatrix} 1 & 0 & 0 & w^7 & 1 & w^{11} & w^{13} & w^4 & w^9 & w^4 & w^{10} & w^{14} \\ 0 & 1 & 0 & 1 & 1 & w^3 & w^7 & w^9 & w^{13} & w^9 & w^8 & \dots & w^2 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \end{bmatrix}$$

Not that  $A_0=1$ ,  $A_{162}=990$ ,  $A_{163}=1155$ ,  $A_{164}=495$ ,  $A_{165}=285$ ,  $A_{166}=435$ ,  $A_{167}=540$ ,  $A_{168}=195$ , and  $S=A_0+A_{162}+A_{163}+A_{164}+A_{165}+A_{166}+A_{167}+A_{168}=4096=16^3$  and the covering radius  $\mu = 166$ .  $\Box$ 

**XI.** (192;14)-arc give linear  $[192,3,178]_{16}$  – *code* define by the generator matrix  $G_{3\times 192}$ .

$$G_{3\times 192} = \begin{bmatrix} 1 & 0 & 0 \ \text{w}^7 & 1 & \text{w}^{11} \ \text{w}^{13} & \text{w}^4 & \text{w}^9 \ \text{w}^4 \ \text{w}^{10} & \text{w}^{10} \\ 0 & 1 & 0 & 1 & 1 & \text{w}^3 \ \text{w}^7 & \text{w}^9 & \text{w}^{13} \ \text{w}^9 \ \text{w}^8 \ \dots \ \text{w}^3 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \end{bmatrix}$$

Not that  $A_0=1$ ,  $A_{178}=915$ ,  $A_{179}=1110$ ,  $A_{180}=780$ ,  $A_{181}=300$ ,  $A_{182}=360$ ,  $A_{183}=435$ ,  $A_{184}=180$ ,  $A_{185}=15$  and  $S=A_0+A_{178}+A_{179}+A_{180}+A_{181}+A_{182}+A_{183}+A_{184}+A_{185}$  $=4096=16^3$  and the covering radius  $\mu = 181$ .  $\Box$ 

**XII.** (213;15)-arc give linear  $[213,3,198]_{16}$  – *code* define by the generator matrix  $G_{3\times 213}$ .

$$G_{3\times213} = \begin{bmatrix} 1 & 0 & 0 & w^7 & 1 & w^{11} & w^{13} & w^4 & w^9 & w^4 & w^{10} & w \\ 0 & 1 & 0 & 1 & 1 & w^3 & w^7 & w^9 & w^{13} & w^9 & w^8 & \dots & w^2 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \end{bmatrix}$$

**XIII.** (234;16)-arc give linear  $[234,3,218]_{16}$  – *code* define by the generator matrix  $G_{3\times 234}$ .

$$G_{3\times134} = \begin{bmatrix} 1 & 0 & 0 & w^7 & 1 & w^{11} & w^{13} & w^4 & w^9 & w^4 & w^{10} & 1 \\ 0 & 1 & 0 & 1 & 1 & w^3 & w^7 & w^9 & w^{13} & w^9 & w^8 & \dots & w^7 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \end{bmatrix}$$

Not that  $A_0=1$ ,  $A_{218}=1290$ ,  $A_{219}=1140$ ,  $A_{220}=945$ ,  $A_{221}=225$ ,  $A_{222}=345$ ,  $A_{223}=135$ ,  $A_{224}=15$ , and  $S=A_0+A_{218}+A_{219}+A_{220}+A_{221}+A_{222}+A_{223}+A_{224}+A_{225}=4096=16^3$  and the covering radius  $\mu = 220$ .  $\Box$ 

XIII. (273;17)-arc give linear  $[273,3,256]_{16}$  – *code* define by the generator matrix  $G_{3\times 273}$ .

$$G_{3\times 273} = \begin{bmatrix} 1 & 0 & 0 & w^7 & 1 & w^{11} & w^{13} & w^4 & w^9 & w^4 & w^{10} & 1 \\ 0 & 1 & 0 & 1 & 1 & w^3 & w^7 & w^9 & w^{13} & w^9 & w^8 & \dots & 0 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \end{bmatrix}$$

Not that  $A_0=1$ ,  $A_{256}=4095$ , and  $S=A_0+A_{256}=4096=16^3$  and the covering radius  $\mu = 259$ .  $\Box$ 

## Conclusion

we constructed a complete (n,r)-arcs in projective plane with respect to r, such that

r = 4,5,..,17 and find the relation between geometrical objects and linear codes.

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